

# Finite element potentials

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## Outline

- 1 Introduction
- 2 The gradient
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## The aim

Beyond a doubt, among the “stars” of vector calculus we have the operators

- **grad**
- **curl**
- **div**

Aim of this talk is to show how to determine in a **constructive** way the **finite element** solutions of  $\mathbf{grad} \psi = \mathbf{H}$ ,  $\mathbf{curl} \mathbf{A} = \mathbf{B}$ ,  $\mathbf{div} \mathbf{v} = G$ .

## First results

Determining the **necessary and sufficient** conditions for assuring that a function defined in a bounded domain  $\Omega \subset \mathbb{R}^3$  is the gradient of a scalar potential, the curl of a vector potential or the divergence of a vector field is one of the **most classical problem** of vector analysis.

The answer is well-known, and shows an interesting interplay of **differential calculus** and **topology** (see, e.g., Cantarella et al. (2002)).

## First results (cont'd)

- a vector field is the **gradient** of a scalar potential if and only if it is **curl free** and its line integral is vanishing on all the closed curves that give a basis of the **first homology group** of  $\overline{\Omega}$ ;
- a vector field is the **curl** of a vector potential if and only if it is **divergence free** and its flux is vanishing across all the closed surfaces that give a basis of the **second homology group** of  $\overline{\Omega}$ , or, equivalently, across (all but one) the connected components of  $\partial\Omega$ ;
- **each** scalar function is the **divergence** of a vector field [just take the gradient of the inverse of the Laplace operator...].

## First results (cont'd)

However, a less clarified situation takes shape when we want to furnish an explicit and efficient procedure for constructing a **finite element** solution.

[**Note:** at this level the divergence case comes back on the table: in fact, the gradient of a (standard) finite element approximate solution of  $\Delta\varphi = G$  has a distributional divergence which is not a function, and therefore this divergence cannot be equal to an assigned finite element.]

We suppose to have:

- for topological reasons
  - a **basis**  $\sigma_n$ ,  $n = 1, \dots, g$ , of the first homology group of  $\overline{\Omega}$ ;
  - a **basis**  $\hat{\sigma}_n$ ,  $n = 1, \dots, g$ , of the first homology group of  $\mathbb{R}^3 \setminus \Omega$ ;
- for the efficiency of the solver
  - a **spanning tree**  $\mathcal{S}_h$  of the graph given by the nodes and the edges of the mesh  $\mathcal{T}_h$ .

[**Note**: a suitable and easy way for constructing  $\sigma_n$  and  $\hat{\sigma}_n$  is presented in Hiptmair and Ostrowski (2002); the determination of a spanning tree is a standard procedure in graph theory.]



## The grad problem

We want to solve  $\mathbf{grad} \psi_h = \mathbf{H}_h$  in the finite element context.

The “right” finite elements are:  $\psi_h$  a **nodal** element,  $\mathbf{H}_h$  an **edge** element.

More precisely, we know that:

- a **Nédélec element** of the lowest order is a vector field in  $H(\mathbf{curl}; \Omega)$  that locally has the form  $\mathbf{a}_K + \mathbf{b}_K \times \mathbf{x}$ ;
- a **curl-free** Nédélec element satisfies  $\mathbf{b}_K = \mathbf{0}$  for each  $K$  (in fact,  $\mathbf{curl}(\mathbf{a}_K + \mathbf{b}_K \times \mathbf{x}) = 2\mathbf{b}_K$ );
- the **gradient** of a (globally continuous) **piecewise-linear** finite element is a vector field in  $H(\mathbf{curl}; \Omega)$  that locally is constant (namely, a curl-free Nédélec element).

## The solution of the grad problem

In other words, for solving  $\mathbf{grad} \psi_h = \mathbf{H}_h$  we have to match two Nédélec edge elements of the lowest order, hence the **line integral** of  $\mathbf{grad} \psi_h$  and  $\mathbf{H}_h$  on each **edge** of the mesh  $\mathcal{T}_h$  has to be the same.

The fundamental theorem of calculus says that

$$\psi_h(v_b) - \psi_h(v_a) = \int_e \mathbf{grad} \psi_h \cdot \boldsymbol{\tau} = \int_e \mathbf{H}_h \cdot \boldsymbol{\tau}$$

for an edge  $e = [v_a, v_b]$ . Hence the linear system associated to  $\mathbf{grad} \psi_h = \mathbf{H}_h$  has exactly **two non-zero values** per row.

## The solution of the grad problem (cont'd)

Starting from the root  $v_*$  of the **spanning tree**  $\mathcal{S}_h$ , where, for the sake of uniqueness, we impose  $\psi_h(v_*) = 0$ , we compute

$$\psi_h(\widehat{v}) = \psi_h(v_*) + \int_{e'} \mathbf{H}_h \cdot \boldsymbol{\tau} \quad (1)$$

for an edge  $e' = [v_*, \widehat{v}] \in \mathcal{S}_h$ ; since  $\mathcal{S}_h$  is a spanning tree, going on in this way we can visit **all** the nodes of  $\mathcal{T}_h$ .

The spanning tree is therefore a tool for selecting the rows for which, using the additional equation  $\psi_h(v_*) = 0$ , one can **eliminate** the unknowns **one after the other**.

We have thus found a nodal element  $\psi_h$  such that its gradient has line integral on all the edges of the **spanning tree** equal to that of  $\mathbf{H}_h$ .

## The solution of the grad problem (cont'd)

What about the edges **not** belonging to the spanning tree?

For each node  $v_i$ ,  $v_i \neq v_*$  let us denote by  $C_{v_i}$  the set of edges in  $\mathcal{S}_h$  joining  $v_*$  to  $v_i$ . Given an edge  $e = [v_a, v_b]$  not belonging to  $\mathcal{S}_h$ , we define the **cycle**  $D_e = C_{v_a} + e - C_{v_b}$ .

Since  $\mathbf{H}_h$  is a gradient (it is **curl-free** and its line integral on all the cycles  $\sigma_n$  **vanishes**), its line integral on  $D_e$  vanishes. Therefore we have

$$\begin{aligned} 0 &= \oint_{D_e} \mathbf{H}_h \cdot d\mathbf{s} = \psi_h(v_a) + \int_e \mathbf{H}_h \cdot \boldsymbol{\tau} - \psi_h(v_b) \\ &= \int_e \mathbf{H}_h \cdot \boldsymbol{\tau} - \int_e \mathbf{grad} \psi_h \cdot \boldsymbol{\tau}. \end{aligned}$$

## The algorithm

The simple **solution algorithm** reads as follows. Let us denote by  $V$  the set of the nodes of the mesh  $\mathcal{T}_h$ , by  $R$  the set of the nodes where the value of  $\psi_h$  is already known, by  $P$  the set of edges  $e \in \mathcal{S}_h$  with exactly one vertex,  $v'(e)$ , not in  $R$ , and by  $E(v)$  the set of edges having  $v$  as a vertex.

### Algorithm

- ①  $R = \{v_*\}$ ,  $P = E(v_*) \cap \mathcal{S}_h$
- ② *while*  $R \neq V$ 
  - ① *pick*  $e \in P$
  - ② *compute*  $\psi_h(v'(e))$  *from* (1)
  - ③ *update*  $P$ :  $P = [P \cup (E(v'(e)) \cap \mathcal{S}_h)] \setminus \{e\}$
  - ④ *update*  $R$ :  $R = R \cup \{v'(e)\}$ .

## The curl problem

We want to solve  $\mathbf{curl} \mathbf{A}_h = \mathbf{B}_h$  in the finite element context.

The “right” finite elements are:  $\mathbf{A}_h$  an **edge** element,  $\mathbf{B}_h$  a **face** element (namely, a Raviart–Thomas element).

More precisely, we know that:

- a **Raviart–Thomas element** of the lowest order is a vector field in  $H(\text{div}; \Omega)$  that locally has the form  $\mathbf{a}_K + b_K \mathbf{x}$ ;
- a **divergence-free** Raviart–Thomas element satisfies  $b_K = 0$  for each  $K$  (in fact,  $\text{div}(\mathbf{a}_K + b_K \mathbf{x}) = 3 b_K$ );
- the **curl** of a (lowest order) **Nédélec element** is a vector field in  $H(\text{div}; \Omega)$  that locally is constant (namely, a divergence-free Raviart–Thomas element).

## The solution of the curl problem

In other words, for solving  $\mathbf{curl} \mathbf{A}_h = \mathbf{B}_h$  we have to match two Raviart–Thomas face elements of the lowest order, hence the **flux** of  $\mathbf{curl} \mathbf{A}_h$  and  $\mathbf{B}_h$  on each **face** of the mesh  $\mathcal{T}_h$  has to be the same.

Since the **Stokes theorem** assures that

$$\int_{e_1} \mathbf{A}_h \cdot \boldsymbol{\tau} + \int_{e_2} \mathbf{A}_h \cdot \boldsymbol{\tau} + \int_{e_3} \mathbf{A}_h \cdot \boldsymbol{\tau} = \int_f \mathbf{curl} \mathbf{A}_h \cdot \boldsymbol{\nu}_f = \int_f \mathbf{B}_h \cdot \boldsymbol{\nu}_f, \quad (2)$$

where  $\partial f = e_1 \cup e_2 \cup e_3$  and  $\boldsymbol{\nu}_f$  is the unit normal vector on  $f$  (with consistent orientation), we deduce that the linear system associated to  $\mathbf{curl} \mathbf{A}_h = \mathbf{B}_h$  has exactly **three non-zero values** for each row.

With respect to the preceding case:

- need to work on the **edges** instead of on the nodes
- more important: **three** unknowns per row instead of two.

## The solution of the curl problem (cont'd)

Having three unknowns per row, in order to devise an **efficient elimination algorithm** it is useful to **fix** the value of other unknowns.

The best situation should occur when number of the new equations agrees with the dimension of the **kernel** of the curl operator.

Since this kernel is given by the **gradients of nodal elements** plus the **basis of the first de Rham cohomology group** of  $\Omega$ , we see that its dimension is equal to  $n_v - 1 + g$  (having denoted by  $n_v$  the number of the nodes in the mesh  $\mathcal{T}_h$ ).



## The solution of the curl problem (cont'd)

Having this in mind, we are led to the problem

$$\begin{aligned}
 \mathbf{curl} \mathbf{A}_h &= \mathbf{B}_h && \text{in } \Omega \\
 \oint_{\sigma_n} \mathbf{A}_h \cdot d\mathbf{s} &= \rho_n && \forall n = 1, \dots, g \\
 \int_{e'} \mathbf{A}_h \cdot \boldsymbol{\tau} &= 0 && \forall e' \in \mathcal{S}_h,
 \end{aligned} \tag{3}$$

for suitable given constants  $\rho_n$  (made precise in the sequel).

[Note that the number of edges  $e'$  in  $\mathcal{S}_h$  is  $n_v - 1$ ; therefore (3)<sub>3</sub> can be seen as a “filter” for gradients. On the other hand, homology and cohomology are in duality, hence (3)<sub>2</sub> can be seen as a “filter” for cohomology fields.]

It is not difficult to prove that there exists a **unique solution** to (3).

## Webb–Forghani algorithm

Webb and Forghani (1989) proposed this **solution algorithm**:

### Algorithm

- 1 *set value 0 to the unknowns corresponding to an edge belonging to the spanning tree*
- 2 *take a face  $f$  for which at least one edge unknown has not yet been assigned*
  - 1 *if exactly one edge unknown is not determined, compute its value from the Stokes relation (2)*
  - 2 *if two or three edge unknowns are not determined, pass to another face*
- 3 *if the iterations stop, use  $\oint_{\sigma_n} \mathbf{A}_h \cdot d\mathbf{s} = \rho_n$  to restart*

[In their case, step 3 was missing, as they considered the case of simple topology (namely,  $g = 0$ ).]

## Webb–Forghani algorithm (cont'd)

The Webb–Forghani algorithm is a simple **elimination procedure** for solving the linear system at hand, and it is quite efficient, as the computational costs is **linearly dependent** on the number of unknowns.

The **weak point** is that:

- it **strongly depends** on the choice of the spanning tree and it can stop without having determined all the edge unknowns (even in simple topological situations!)

(see Dłotko and Specogna (2010)).

## An explicit formula for the vector potential

- **Cure:** devise an **explicit formula** for the solution to (3).

(We are able to do that if  $\mathbf{B}_h \cdot \boldsymbol{\nu} = 0$  on  $\partial\Omega$ , a quite natural condition in the most interesting physical situations.)

The idea is the following. Define the **Biot–Savart** field

$$\mathbf{H}^{BS}(\mathbf{x}) = \frac{1}{4\pi} \int_{\Omega} \mathbf{B}_h(\mathbf{y}) \times \frac{\mathbf{x} - \mathbf{y}}{|\mathbf{x} - \mathbf{y}|^3} d\mathbf{y},$$

and set  $\rho_n = \oint_{\sigma_n} \mathbf{H}^{BS} \cdot d\mathbf{s}$  in (3).

One has  $\mathbf{curl} \mathbf{H}^{BS} = \mathbf{B}_h$  in  $\Omega$  (here the condition  $\mathbf{B}_h \cdot \boldsymbol{\nu} = 0$  on  $\partial\Omega$  has played a role). Hence the Nédélec interpolant  $\Pi^{N_h} \mathbf{H}^{BS}$  satisfies (3)<sub>1</sub> and (3)<sub>2</sub>.

To find the solution to (3), we can correct  $\Pi^{N_h} \mathbf{H}^{BS}$  by a gradient, namely, construct the nodal element whose gradient has the **same line integral** of  $\mathbf{H}^{BS}$  on the edges of the spanning tree  $\mathcal{S}_h$ .

## An explicit formula for the vector potential (cont'd)

Being  $v_*$  the root of the spanning tree, in the nodes of the mesh define the finite element  $\phi_h$  as  $\phi_h(v_*) = 0$  and

$$\phi_h(v_b) = \phi_h(v_a) + \int_{e'} \mathbf{H}^{BS} \cdot \boldsymbol{\tau} \quad \forall e' = [v_a, v_b] \in \mathbf{S}_h.$$

The Nédélec finite element  $\mathbf{A}_h = \Pi^{N_h} \mathbf{H}^{BS} - \mathbf{grad} \phi_h$  is the **solution** to (3).

To express its **degrees of freedom**, we proceed as follows. For each edge  $e \notin \mathcal{S}_h$ , we define the **cycle**  $D_e$  as before (the edges from the root of the spanning tree to the first vertex of  $e$ , the edge  $e$ , the edges from the second vertex of  $e$  to the root of the spanning tree).

## An explicit formula for the vector potential (cont'd)

The cycle  $D_e$  is constituted by edges **all belonging to the spanning tree** (except  $e$ ): hence we have

$$\begin{aligned}
 \int_e \mathbf{A}_h \cdot \boldsymbol{\tau} &= \int_e (\Pi^{N_h} \mathbf{H}^{BS} - \mathbf{grad} \phi_h) \cdot \boldsymbol{\tau} \\
 &= \int_e \mathbf{H}^{BS} \cdot \boldsymbol{\tau} - [\phi_h(v_b) - \phi_h(v_a)] \\
 &= \int_e \mathbf{H}^{BS} \cdot \boldsymbol{\tau} - \left[ \int_{C_{v_b}} \mathbf{H}^{BS} \cdot \boldsymbol{\tau} - \int_{C_{v_a}} \mathbf{H}^{BS} \cdot \boldsymbol{\tau} \right] \quad (4) \\
 &= \oint_{D_e} \mathbf{H}^{BS} \cdot d\mathbf{s} \\
 &= \frac{1}{4\pi} \oint_{D_e} \left( \int_{\Omega} \mathbf{B}_h(\mathbf{y}) \times \frac{\mathbf{x}-\mathbf{y}}{|\mathbf{x}-\mathbf{y}|^3} d\mathbf{y} \right) \cdot d\mathbf{s}(\mathbf{x}).
 \end{aligned}$$

Using (4), we can always **restart** the Webb–Forghani algorithm.

## A basis of the first de Rham cohomology group

This algorithm permits to solve also the problem

$$\begin{aligned} \mathbf{curl} \mathbf{A}_h &= \mathbf{0} && \text{in } \Omega \\ \oint_{\sigma_n} \mathbf{A}_h \cdot d\mathbf{s} &= \kappa_n && \forall n = 1, \dots, g \\ \int_{e'} \mathbf{A}_h \cdot \boldsymbol{\tau} &= 0 && \forall e' \in \mathcal{S}_h, \end{aligned} \quad (5)$$

for any choice of the constants  $\kappa_n$ . In particular, taking  $\kappa_n$  equal to the entries  $Q_{n,j}$  of a non-singular  $g \times g$  square matrix  $Q$ , we find a basis  $\mathbf{T}^{(j)}$  of the **first de Rham cohomology group**.

In fact, an **explicit formula** is available also in this case: the choice of  $Q_{n,j} = \ell_\kappa(\sigma_n, \hat{\sigma}_j)$  gives

$$\int_e \mathbf{T}^{(j)} \cdot \boldsymbol{\tau} = \ell_\kappa(D_e, \hat{\sigma}_j),$$

where  $\ell_\kappa$  is the **linking number** (and  $\hat{\sigma}_j$  has been chosen inside  $\mathbb{R}^3 \setminus \overline{\Omega}$ , namely, not intersecting  $\partial\Omega$ ).

## The linking number

The **linking number** between  $\hat{\sigma}_j$  and another **disjoint** cycle  $\sigma$  is given by:

$$l_{\mathcal{K}}(\sigma, \hat{\sigma}_j) = \frac{1}{4\pi} \oint_{\sigma} \left( \oint_{\hat{\sigma}_j} \frac{\mathbf{y} - \mathbf{x}}{|\mathbf{y} - \mathbf{x}|^3} \times d\mathbf{s}_y \right) \cdot d\mathbf{s}_x.$$

- The linking number (introduced by Gauss...) is an **integer** that represents the number of times that each cycle **winds** around the other.



## The divergence problem

We want to solve  $\operatorname{div} \mathbf{v}_h = G_h$  in the finite element context.

The “right” finite elements are:  $\mathbf{v}_h$  a **face** element,  $G_h$  a **nodal** element.

More precisely, we know that:

- a **nodal element** of the lowest order is a function in  $L^2(\Omega)$  that is locally constant;
- the **divergence** of a (lowest order) **Raviart–Thomas** finite element is a function in  $L^2(\Omega)$  that is locally constant.

## The solution of the divergence problem

In other words, for solving  $\operatorname{div} \mathbf{v}_h = G_h$  we have to match two piecewise-constant elements, hence the **integral** of  $\operatorname{div} \mathbf{v}_h$  and  $G_h$  on each **element** of the mesh  $\mathcal{T}_h$  has to be the same.

The **Gauss theorem** says that

$$\begin{aligned} \int_{f_1} \mathbf{v}_h \cdot \boldsymbol{\nu}_K + \int_{f_2} \mathbf{v}_h \cdot \boldsymbol{\nu}_K + \int_{f_3} \mathbf{v}_h \cdot \boldsymbol{\nu}_K + \int_{f_4} \mathbf{v}_h \cdot \boldsymbol{\nu}_K \\ = \int_K \operatorname{div} \mathbf{v}_h = \int_K G_h, \end{aligned} \quad (6)$$

where  $\partial K = f_1 \cup f_2 \cup f_3 \cup f_4$  and  $\boldsymbol{\nu}_K$  is the unit outward normal vector on  $\partial K$ .

Hence the linear system associated to  $\operatorname{div} \mathbf{v}_h = G_h$  has exactly **four unknowns** per row.

## The solution of the divergence problem (cont'd)

In order to **reduce** the dimension of the system, we want to **fix** the value of some unknowns. Similarly to what done before we start by analyzing the **dimension** of the **kernel** of the divergence operator.

This kernel is given by the curl of the Nédélec elements plus the basis of the **second** de Rham cohomology group of  $\Omega$ .

If we denote by  $(\partial\Omega)_0, \dots, (\partial\Omega)_p$  the **connected components** of  $\partial\Omega$ , we know that the dimension of the second de Rham cohomology group of  $\Omega$  is equal to  $p$ .

## The solution of the divergence problem (cont'd)

On the other hand, it is easy to check that the **dimension** of the space of the curl of the Nédélec elements is equal to the number of the edges minus the dimension of the kernel of the curl operator. Hence, it is equal to  $n_e - n_v + 1 - g$ .

By the **Euler–Poincaré formula** we have

$$n_v - n_e + n_f - n_t = 1 - g + p,$$

hence the dimension of the space of the curl can be rewritten as  $n_f - n_t - p$ .

In conclusion, besides the **topological** conditions

$$\int_{(\partial\Omega)_r} \mathbf{v}_h \cdot \boldsymbol{\nu} = c_r, \quad r = 1, \dots, p,$$

that are a **filter for the cohomology fields**, we could add  $n_f - n_t - p$  equations.

## A dual graph

To do that, let us note that an internal face **connects** two tetrahedra, while a boundary face **connects** a tetrahedron and a connected component of  $\partial\Omega$ .

We can therefore consider the following (connected) **dual graph**  $\mathcal{G}_h$ : the dual **vertices** are  $W = T \cup \Sigma$ , where the elements of  $T$  are the tetrahedra of the mesh and the elements of  $\Sigma$  are the  $p + 1$  connected components of  $\partial\Omega$ ; the set of dual **arcs** is  $F$ , the set of the faces of the mesh.

## A dual graph (cont'd)

The **number** of dual vertices is equal to  $n_t + p + 1$ , hence a **spanning tree**  $\mathcal{M}_h$  of  $\mathcal{G}_h$  has  $n_t + p$  dual arcs (and consequently its **cotree** has  $n_f - n_t - p$  dual arcs).

Therefore the linear system

$$\begin{aligned} \operatorname{div} \mathbf{v}_h &= G_h && \text{in } \Omega \\ \int_{(\partial\Omega)_r} \mathbf{v}_h \cdot \boldsymbol{\nu} &= c_r && \forall r = 1, \dots, p \\ \int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f &= 0 && \forall f \notin \mathcal{M}_h \end{aligned} \quad (7)$$

is a **square linear system** of  $n_f$  equations and unknowns.

## Well-posedness of (7)

Now we show that it has a **unique solution**.

The procedure is **constructive**, similar in some sense to the elimination procedure used for the grad problem but now going along the **dual spanning tree**, starting from the **leaves**. (Let us recall that the leaves of a spanning tree  $\mathcal{M}_h$  are the vertices of  $W$  that have **only one arc** of  $\mathcal{M}_h$  incident to them.)

Remembering that we have imposed  $\int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f = 0$  if  $f \notin \mathcal{M}_h$ , we can reduce the problem to the faces  $f \in \mathcal{M}_h$ .

## Well-posedness of (7) (cont'd)

Given  $w \in W$  (a **tetrahedron** or a **connected component** of  $\partial\Omega$ ), let us set  $\mathcal{F}(w) = \{f \in F : f \subset w\}$ ; the elements of this set are faces of the primal mesh, therefore dual arcs in the dual mesh.

The **leaves** of the spanning tree  $\mathcal{M}_h$  are the vertices  $w \in W$  such that  $\mathcal{F}(w) \cap \mathcal{M}_h$  reduces to exactly one dual arc (namely, to a face).



## Well-posedness of (7) (cont'd)

If  $w$  is a **leave** of  $\mathcal{M}_h$  and  $f(w)$  is the unique dual arc in  $\mathcal{M}_h$  incident to  $w$ , we can easily compute the **degree of freedom** corresponding to  $f(w)$ , as we know that  $\int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f = 0$  for all  $f \notin \mathcal{M}_h$ .

In fact we have

$$\int_{f(w)} \mathbf{v}_h \cdot \boldsymbol{\nu}_f = \begin{cases} \int_{\partial w} \mathbf{v}_h \cdot \boldsymbol{\nu}_f = \int_w G_h & \text{if } w \in T \\ \int_{(\partial\Omega)_r} \mathbf{v}_h \cdot \boldsymbol{\nu} = c_r & \text{if } w = (\partial\Omega)_r, \\ & r = 1, \dots, p \\ \int_{(\partial\Omega)_0} \mathbf{v}_h \cdot \boldsymbol{\nu} = \int_{\Omega} G_h - \sum_{r=1}^p c_r & \text{if } w = (\partial\Omega)_0, \end{cases}$$

having used the **Gauss theorem** in the first and the third line.

## Well-posedness of (7) (cont'd)

Hence it is clear that, if  $\mathbf{v}_h$  is a Raviart–Thomas element with  $\operatorname{div} \mathbf{v}_h = 0$ ,  $\int_{(\partial\Omega)_r} \mathbf{v}_h \cdot \boldsymbol{\nu} = 0$  for all  $r = 1, \dots, p$ , and  $\int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f = 0$  for all  $f \notin \mathcal{M}_h$ , then  $\int_{f(w)} \mathbf{v}_h \cdot \boldsymbol{\nu}_f = 0$  for the faces  $f(w)$  associated to all the leaves  $w \in \mathcal{M}_h$ .

We can **iterate** this argument: if we remove from the spanning tree  $\mathcal{M}_h$  a leaf and its corresponding incident arc, the remaining graph is **still** a tree.

**Repeating** the previous procedure, we can easily compute the degrees of freedom corresponding to the faces incident to the leaves of this new tree, finding that they are vanishing.

After a finite number of steps the remaining tree reduces to just one vertex, and the result is that  $\int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f = 0$  for all  $f \in F$ . This proves that (7) has a **unique solution**.

## The solution algorithm

We can also furnish an **explicit way** for computing the values of the degrees of freedom.

In fact, let us denote by  $U$  the set of vertices  $w$  and by  $N$  the set of arcs  $f$  of the **reduced** dual graph obtained at a step of the previous procedure (at the initial step, the set of arcs of the spanning tree is denoted by  $M$ ).

Then:

- $N$  is the set of faces where the degree of freedom  $\int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f$  is still **unknown**;
- if  $w$  is a leave of  $(U, N)$ , there exists **exactly one face**  $f(w)$  incident to it and belonging to  $N$ .

## The solution algorithm (cont'd)

Then, if  $w$  is a leave of  $(U, N)$ , we have

$$\int_{f(w)} \mathbf{v}_h \cdot \boldsymbol{\nu}_f = A_w - \sum_{f \in \mathcal{F}(w) \setminus f(w)} \int_f \mathbf{v}_h \cdot \boldsymbol{\nu}_f, \quad (8)$$

where

$$A_w = \begin{cases} \int_w G_h & \text{if } w \in T \\ c_r & \text{if } w = (\partial\Omega)_r, r = 1, \dots, p \\ \int_\Omega G_h - \sum_{r=1}^p c_r & \text{if } w = (\partial\Omega)_0. \end{cases}$$

## The solution algorithm (cont'd)

This can be rephrased as an **elimination algorithm** for the computation of  $\mathbf{v}_h$ .

### Algorithm

- ①  $N = M, U = W$
- ② *while*  $N \neq \emptyset$ 
  - ① *pick a leaf*  $w$  *of the tree*  $(U, N)$
  - ② *compute*  $\int_{f(w)} \mathbf{v}_h \cdot \boldsymbol{\nu}_f$  *from* (8)
  - ③ *update*  $U$ :  $U = U \setminus \{w\}$
  - ④ *update*  $N$ :  $N = N \setminus \{f(w)\}$ .

Notice that at any step of the algorithm  $(U, N)$  is a tree, so while  $N \neq \emptyset$  a leaf  $w$  of  $(U, N)$  **always exists**.

## A basis of the second de Rham cohomology group

- It is worth noting that the set of vector fields  $\mathbf{W}^{(s)}$ ,  $s = 1, \dots, p$ , solutions to problem (7) with  $G_h = 0$  and  $c_r = \delta_{r,s}$ ,  $r = 1, \dots, p$ , is a **basis** of the **second de Rham cohomology group** of  $\Omega$ .

## Stability

A natural question is: how to construct **stable** finite element potentials?

(This means potentials whose natural norms can be estimated in terms of the norms of the data, **uniformly** with respect to the mesh size  $h$ .)

## Stability (cont'd)

[However, let us note that, very often, the construction of finite element potentials is a **preliminary step** in the procedure aiming at solving a partial differential equation.

In this respect, the solution  $u_h$  will be written as  $u_h = U_h + W_h$ ,  $W_h$  being the finite element **potential** and  $U_h$  the solution of an **auxiliary problem** in which  $W_h$  contributes at the **right hand side**.

In this situation, what is **interesting** is the stability of the solution  $u_h$ , and not that of  $W_h$  and  $U_h$ ; in many cases, an unstable  $W_h$  produces an unstable  $U_h$  but a **stable**  $u_h$ .]



## Stability (cont'd)

- **Grad** problem: **project** the scalar potential  $\psi_h$  on the space orthogonal to the **constants**, namely, take

$$\psi_h^* = \psi_h - \frac{1}{\text{meas}(\Omega)} \int_{\Omega} \psi_h.$$

This is a **stable** scalar potential, satisfying

$$\|\psi_h^*\|_{H^1(\Omega)} \leq K_G \|\mathbf{H}_h\|_{(L^2(\Omega))^3}.$$

## Stability (cont'd)

- **Curl** problem: **project** the vector potential  $\mathbf{A}_h$  (solution to (3) with  $\rho_n = 0$ ) on the space orthogonal to the **gradients** of the piecewise-linear nodal finite elements, namely, take

$$\mathbf{A}_h^* = \mathbf{A}_h - \mathbf{grad} \Phi_h,$$

where  $\Phi_h$  is the solution to

$$\int_{\Omega} \mathbf{grad} \Phi_h \cdot \mathbf{grad} \eta_h = \int_{\Omega} \mathbf{A}_h \cdot \mathbf{grad} \eta_h \quad \forall \eta_h.$$

This is a **stable** vector potential, satisfying

$$\|\mathbf{A}_h^*\|_{H(\mathbf{curl}; \Omega)} \leq K_C \|\mathbf{B}_h\|_{(L^2(\Omega))^3}.$$

## Stability (cont'd)

- **Divergence** problem: **project** the potential  $\mathbf{v}_h$  (solution to (7) with  $c_r = 0$ ) on the space orthogonal to the **curls** of the lowest order Nédélec finite elements, namely, take

$$\mathbf{v}_h^* = \mathbf{v}_h - \mathbf{curl} \mathbf{q}_h,$$

where  $\mathbf{q}_h$  is the solution to

$$\begin{aligned} \int_{\Omega} \mathbf{curl} \mathbf{q}_h \cdot \mathbf{curl} \mathbf{p}_h &= \int_{\Omega} \mathbf{v}_h \cdot \mathbf{curl} \mathbf{p}_h & \forall \mathbf{p}_h \\ \int_{\Omega} \mathbf{q}_h \cdot \mathbf{grad} \eta_h &= 0 & \forall \eta_h \\ \oint_{\sigma_n} \mathbf{q}_h \cdot d\mathbf{s} &= 0 & \forall n = 1, \dots, g. \end{aligned}$$

This is a **stable** potential, satisfying

$$\|\mathbf{v}_h^*\|_{H(\text{div}; \Omega)} \leq K_D \|G_h\|_{L^2(\Omega)}.$$

## Stability (cont'd)

All these stability results are based on the fact that a **Poincaré-like inequality** is valid for functions orthogonal to constants, gradients and curls, respectively:

$$\begin{aligned} \|\psi_h^*\|_{L^2(\Omega)} &\leq K_G^* \|\mathbf{grad} \psi_h^*\|_{(L^2(\Omega))^3} \\ \|\mathbf{A}_h^*\|_{(L^2(\Omega))^3} &\leq K_C^* \|\mathbf{curl} \mathbf{A}_h^*\|_{(L^2(\Omega))^3} \\ \|\mathbf{v}_h^*\|_{(L^2(\Omega))^3} &\leq K_D^* \|\mathbf{div} \mathbf{v}_h^*\|_{L^2(\Omega)}. \end{aligned}$$

[To be precise, here the additional **topological** conditions  $\oint_{\sigma_n} \mathbf{A}_h^* \cdot d\mathbf{s} = 0$  for  $n = 1, \dots, g$ , and  $\int_{(\partial\Omega)_r} \mathbf{v}_h^* \cdot \boldsymbol{\nu} = c_r$  for  $r = 1, \dots, p$ , also play a role.]

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